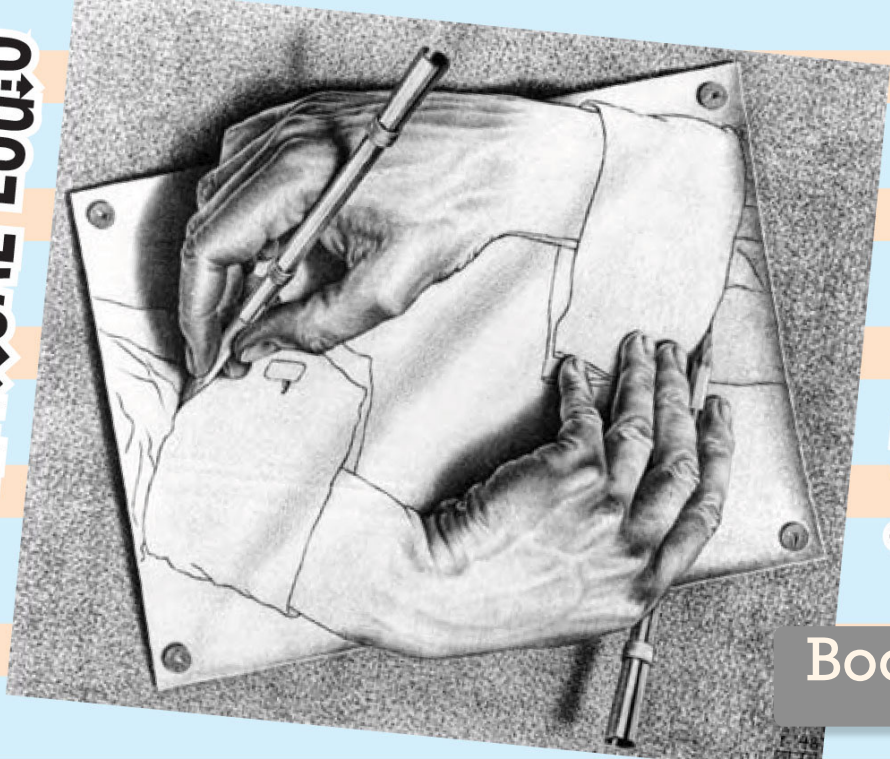


MATHEMATICAL LOGIC



& SET THEORY

Book 1

Group Theory: an example of a first-order axiomatic system

An informal proof in group theory

Theorem If G is a (multiplicative) group of exponent 2, then G is abelian.

(G has exponent n if $g^n = 1$ for all $g \in G$.)

(Informal) proof: Let $a, b \in G$. Since $abab = (ab)^2 = 1$, multiplying on the left by "a" and on the right by "b" gives $aababb = a1b$, i.e. $ba = ab$. \square

Axioms of Group Theory:

ID: $(\forall x) ((x * 1 = x) \wedge (1 * x = x))$

ASSOC: $(\forall x)(\forall y)(\forall z) ((x * y) * z = x * (y * z))$

INV: $(\forall x)(\exists y) ((x * y = 1) \wedge (y * x = 1))$

i.e. $\mu(\mu(x,y),z) = \mu(x,\mu(y,z))$

Start with names for variables x, y, z, \dots (symbols)
Special symbols for first order logic: \exists, \forall , parentheses, \neg, \rightarrow, \dots

Symbols for constants: $1, \dots$

Symbols for functions: $*$, ... $x * y$ means $\mu(x, y)$

Symbols for relations: $=$

We happen to know some groups including C_n (cyclic group of order n), S_n (symmetric group of degree n), ...

GROUPS = $\{ID, ASSOC, INV\} = \{(\forall x)((x * 1) = \dots, \dots, \dots)\}$ (the set consisting of our three axioms of group theory)

S_5 is a group, i.e. $S_5 \models$ GROUPS (S_5 is a model of GROUPS)

ABEL: $(\forall x)(\forall y) (x * y = y * x)$

ABEL-GPS = GROUPS \cup {ABEL}. S_5 is a non-abelian group; $S_5 \not\models$ ABEL; $S_5 \not\models$ ABEL-GPS.

A structure has an underlying set of elements, together with an interpretation of all the symbols for constants, functions, and relations.

How do we rewrite our informal proof (above) as a formal proof in first order logic?

$\Sigma = \text{GROUPS} \cup \{\text{EXP2}\}$ where $\text{EXP2}: (\forall x)(x*x=1)$

ABEL is a theorem in the theory of groups of exponent 2, i.e. $\Sigma \vdash \text{ABEL}$.

A theorem is a sequence of steps $\Sigma \vdash \square$ in which every step follows from previous steps by a statement in Σ , or an axiom of first order logic, or a rule of inference.

$\Sigma \vdash \square$
 $\Sigma \vdash \square$
 $\Sigma \vdash \square$
 \vdots
 $\Sigma \vdash \square$

This is a formal (symbolic) proof!

An outline of a formal proof: $\Sigma \vdash \text{EXP2}$ since $\text{EXP2} \in \Sigma$

$\Sigma \vdash (\text{EXP2} \rightarrow (\forall a)(a*a=1))$ (A4) p.86

$\Sigma \vdash (\forall a)(a*a=1)$ Modus Ponens (R1) p.86

\vdots
 $\Sigma \vdash (\forall b)(b*b=1)$

\vdots
 $\Sigma \vdash (\forall a)(\forall b)((a*b)*(a*b)=1)$

\vdots
 $\Sigma \vdash (\forall a)(\forall b)((a*(a*b)*(a*b))=a*1)$

\vdots
 $\Sigma \vdash (\forall a)(\forall b)(a*b=b*a)$

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ORD3: $(\exists x)(\exists y)(\exists z)[(\forall q)((q=x) \vee (q=y) \vee (q=z)) \wedge (\overset{x \neq y}{\neg(x=y)}) \wedge (\overset{x \neq z}{\neg(x=z)}) \wedge (\overset{y \neq z}{\neg(y=z)})]$

"there are at most three elements"

"there are at least 3 elements"

ABEL is independent of GROUPS (you cannot either prove or disprove that a general group is abelian). GROUPS $\not\vdash$ ABEL and GROUPS $\not\vdash \neg$ ABEL. This is because $C_3 \models \text{GROUPS}$ but $C_3 \not\models \text{ABEL}$ and $S_3 \models \text{GROUPS}$ but $S_3 \not\models \text{ABEL}$.

In an arbitrary first-order theory, with axioms Σ , a statement θ is independent of Σ if

$\Sigma \not\vdash \theta$ and $\Sigma \not\vdash \neg\theta$:

Soundness Theorem: If $\Sigma \vdash \theta$ then θ holds in every model of Σ i.e. $M \models \theta$ whenever $M \models \Sigma$.

Completeness Theorem: Converse holds: If θ holds in every model of Σ , then it is provable from Σ i.e. if $M \models \theta$ whenever $M \models \Sigma$, then $\Sigma \vdash \theta$.

Assume Σ is consistent

So: θ is independent of Σ iff there are models of Σ in which θ holds, and models of θ in which θ fails.

Σ is consistent if we cannot prove a contradiction from Σ , i.e. $\Sigma \not\vdash (\theta \wedge \neg\theta)$ for some θ .

Equivalently, Σ is consistent iff it has a model.

Eq. ABEL is independent of GROUPS.

ORDS

GROUPS is consistent.

GROUPS \cup {ORDS} is consistent since it has a model. In fact it has a unique model up to isomorphism: the cyclic group C_3 of order 3. The group C_3 (or its theory) is categorical.
GROUPS is not categorical. (There are models, but not a unique model.)

An alternative to MV: $(\forall x)(\exists y)((x*y=1) \wedge (y*x=1))$ is to add a function symbol $\iota(\cdot)$ to the language
namely $(\forall x)((x*\iota(x)=1) \wedge (\iota(x)*x=1))$
We already have a binary function symbol $\mu(\cdot, \cdot)$, $\mu(x,y) = x*y$

A theorem of Σ is a statement that can be proved from Σ . A proof is a sequence of statements such....
The theory of Σ is $Th(\Sigma) = \{ \text{statements provable from } \Sigma \} = \{ \text{theorems of } \Sigma \}$.