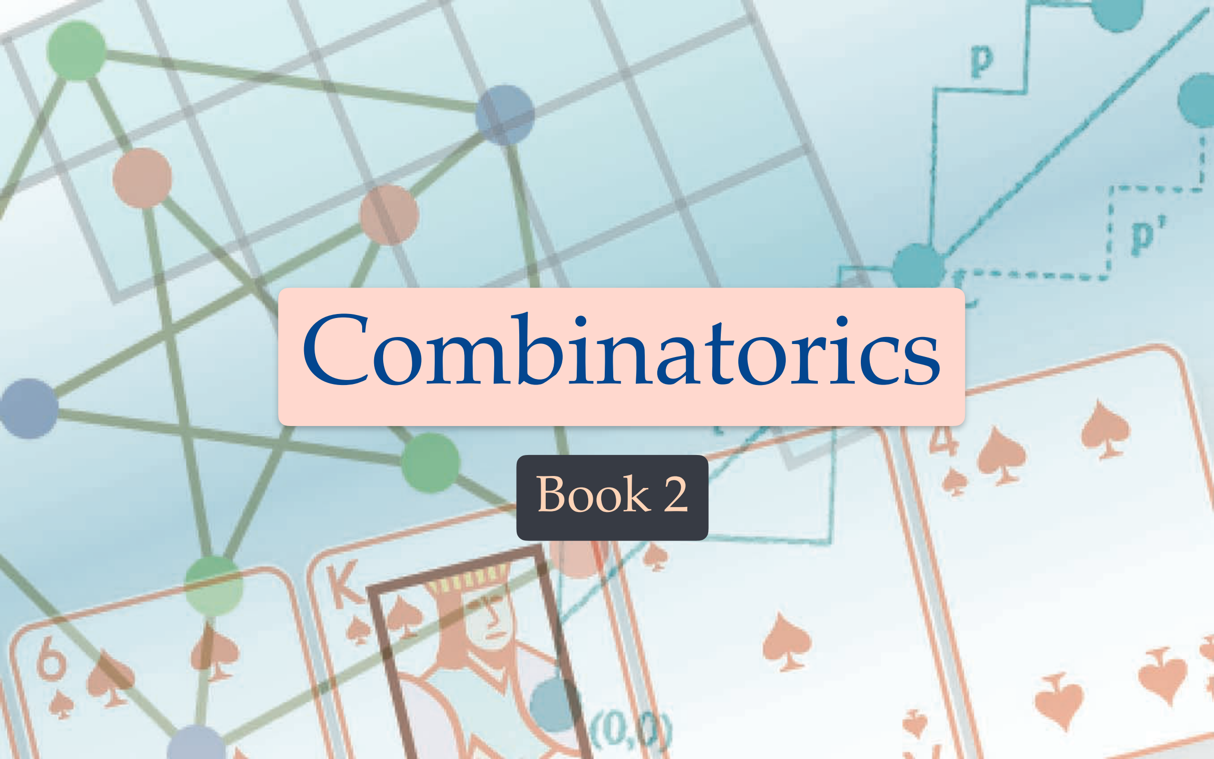
Combinatorics

Book 2



The graph $\Gamma \setminus v$ (formed by removing v and its edges from Γ) has one fewer vertex, so it can be properly colored using at most 6 colors. And since v has at most 5 neighbors in $\Gamma \setminus v$, there is a color left over which can be used to color vertex v. This gives a proper coloring of Γ using at most 6 colors (a contradiction ...)

We will improve this to show that actually 5 colors suffice to properly color every planar graph.

$\psi \ \phi \ \chi \ \pi$

A b c ... x X χ

Greek "chi"

Given a graph Γ, the chromatic number of Γ, denoted $\chi(\Gamma)$, is the smallest number of colors we can use to properly color the vertices of Γ. A proper coloring of the vertices of Γ is a coloring of the vertices such that no edge has both endpoints of the same color.

The theorem of Appel and Haken (1976) is that every planar graph Γ has $\chi(\Gamma) \leq 4$.

Note that $\chi(K_n) = n$. Here K_n is the complete graph of order n.

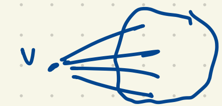
A graph Γ has $\chi(\Gamma) = 1$ iff it has vertices but no edges.

A graph Γ has $\chi(\Gamma) \leq 2$ iff Γ is bipartite iff Γ has no circuits of odd length.

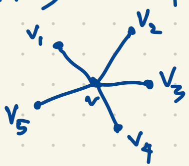
Computing $\chi(\Gamma)$ is hard in general.

Theorem If Γ is a finite planar graph then $\chi(\Gamma) \leq 5$. Proof due to Heawood.

Proof If the theorem fails then there is a smallest counterexample Γ with n vertices (so Γ is planar and every planar graph of order $n-1$ has chromatic number ≤ 5 while $\chi(\Gamma) \ge 6$). We seek a contradiction. Γ has a vertex v of degree ≤ 5. In fact $\deg v = 5$. (If $\deg v \le 4$ then $\chi(\Gamma) \le 5$, a contradiction.) Let Γ' be the graph obtained from Γ by deleting v and its five edges,



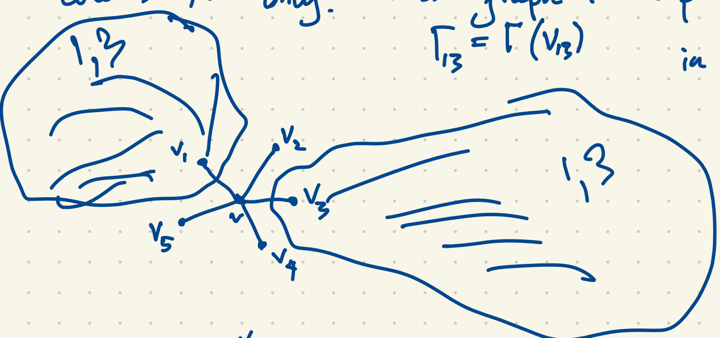
$n-1$ vertices can be properly colored using at most five colors.



so $\chi(\Gamma') \le 5$. Say v_i has color i $(i = 1, 2, \cdots, 5)$.

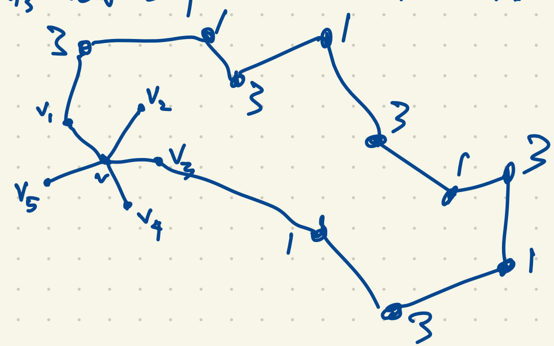
Consider the vertices $V_{13} \subset$ {vertices of Γ} having colors 1, 3 only. This graph is bipartite. I can assume v_1 is joined to v_3 in Γ_{13} (otherwise

$\Gamma_{13} = \Gamma(V_{13})$



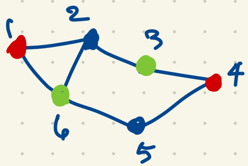
in part of Γ_{13}, reverse colors 1, 3 so that v_3 gets color 1. Then we are free to color v using color 3 since its neighbors are color 1, 2, 1, 4, 5).

Otherwise Γ_{13} has a path from v_1 to v_3.

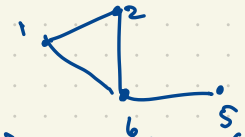


Similarly there is a path from v_2 to v_4 using only vertices of colors 2 and 4. Contradiction! □

Given a graph Γ, a <u>subgraph</u> of Γ is formed by taking a subset of the edges of Γ together with all their vertices. An <u>induced subgraph</u> of Γ is formed by taking a subset of the vertices of Γ together with all their edges in Γ

$\Gamma =$ 

eg.  is a subgraph of Γ. (not an induced subgraph of Γ)

eg.  is an induced subgraph of Γ.

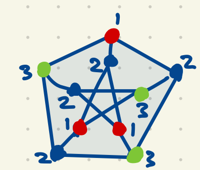
An induced subgraph of Γ is a subgraph of Γ, but not conversely.

A <u>k-clique</u> in Γ is a complete subgraph of Γ, i.e. a subset of the vertices, any two of which are joined.

In Γ above, $\{1, 2, 6\}$ is a clique (in fact a 3-clique). The <u>clique number</u> of Γ, denoted $\omega(\Gamma)$, is the size of the largest clique in Γ. It is hard to compute $\omega(\Gamma)$.

w vs. ω
Roman Greek

<u>Theorem</u> For every graph Γ, $\chi(\Gamma) \geq \omega(K)$.



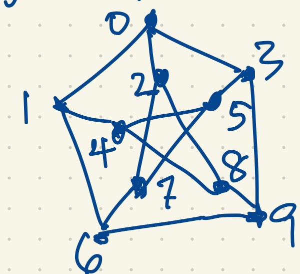
$\chi(P) = 3$ for the Petersen graph P.

Warning: this not equality! For the Petersen graph P, $\omega(P) = 2$.

Proof: The vertices in a clique of size $\omega(\Gamma)$ require $\omega(\Gamma)$ different colors.

Dual to the clique number $\omega(\Gamma)$ we have the coclique number $\alpha(\Gamma)$ which is the maximum number of vertices in Γ, no two of which are joined. (This is $\alpha(\Gamma) = \omega(\bar{\Gamma})$ where $\bar{\Gamma}$ is the complementary graph). Cocliques are also called independent sets of vertices.

Eg. $\alpha(P) = 4$. $|\chi(P)| \geq \frac{10}{4} = 2.5 \Rightarrow \chi(P) \geq 3$.



$\{1,2,3\}$ is a coclique which is not contained in any larger coclique; it is a <u>maximal</u> coclique.

A <u>maximum</u> coclique (i.e. a coclique of maximum size) is $\{1,3,7,8\}$.

This is maximum size because P has vertex set $\{0,3,9,6,1\} \cup \{2,7,5,4,8\}$ as a union of two 5-cycles (circuits of length 5). Any set of size at least 5 vertices has either 3 on the inner 5-cycle $\{2,7,5,4,8\}$ or 3 vertices on the outer 5-cycle $\{0,3,9,6,1\}$. In either case there is an edge in that 5-cycle joining two vertices we have chosen.

<u>Theorem</u> $\chi(\Gamma) \geq \dfrac{|V|}{\alpha(\Gamma)}$ where $|V|$ = the number of vertices = the order of Γ.

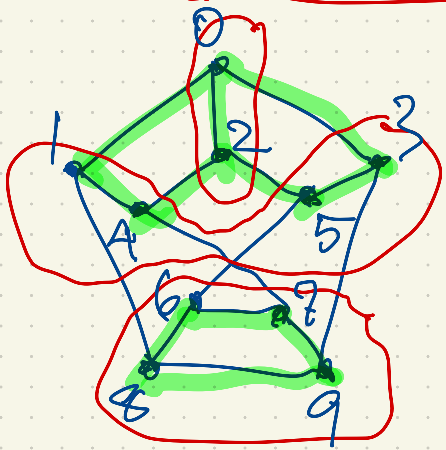
<u>Proof</u> Let $k = \chi(\Gamma)$. Properly color the vertices $1, 2, \cdots, k$ and let V_i be the subset of vertices colored i, for $i = 1, 2, \cdots, k$. This gives a partition $V = V_1 \sqcup V_2 \sqcup \cdots \sqcup V_k$ ($V_i \cup V_j$ = union of V_i and V_j; $V_i \sqcup V_j$ = disjoint union of V_i and V_j). Each V_i is a coclique so $|V_i| \leq \alpha(\Gamma)$ so $|V| = |V_1| + |V_2| + \cdots + |V_k| \leq \underbrace{\alpha(\Gamma) + \alpha(\Gamma) + \cdots + \alpha(\Gamma)}_{k} = k\,\alpha(\Gamma)$ so $\chi(\Gamma) = k \geq \dfrac{|V|}{\alpha(\Gamma)}$. $\square$

March

M	W	F
	1	3
6	$\boxed{8}$	10
13	15	17

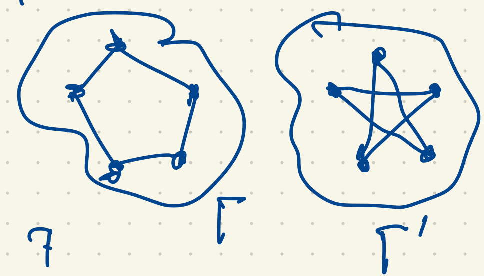
Spring Break

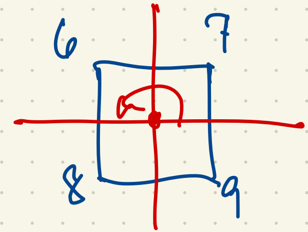
Test 1: Wed Mar 8



$= \Gamma$

You can use nauty to test isomorphism between two graphs.





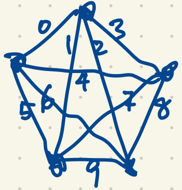
Using nauty Software

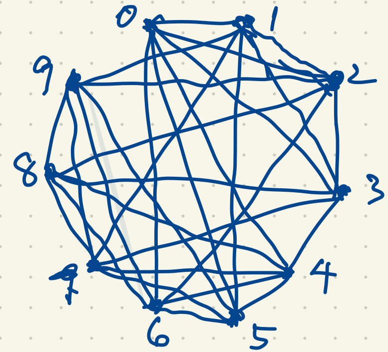
$G = \operatorname{Aut} \Gamma$

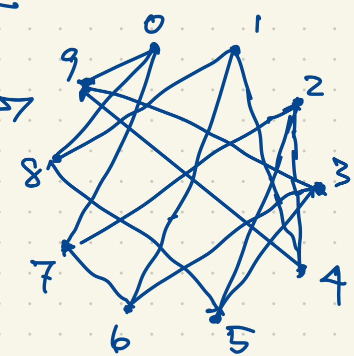
$G = \langle (1,3)(4,5)(6,7)(8,9),\ (0,2)(1,4)(3,5)(6,9)(7,8) \rangle$

$|G| = 4$

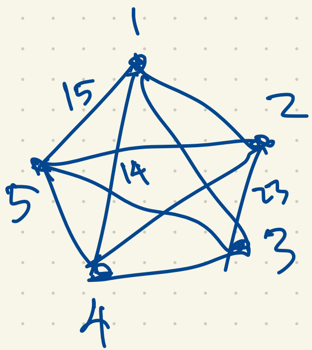
G has 3 orbits on the vertices: $\{0,2\}$, $\{1,3,4,5\}$, $\{6,7,8,9\}$

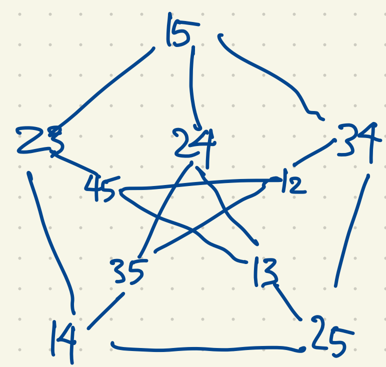
$K_5 =$ 

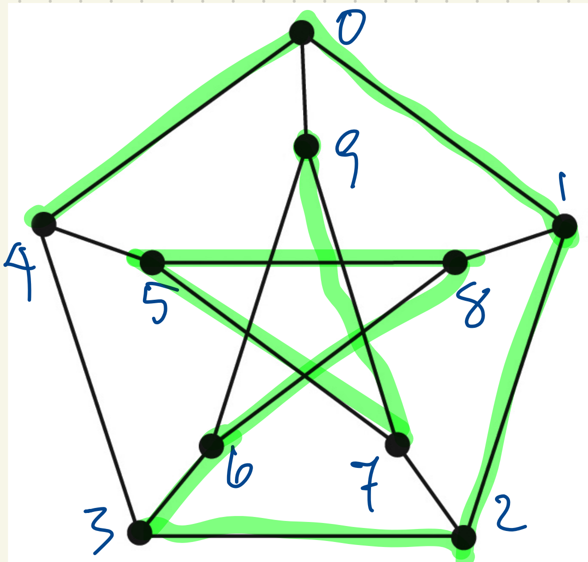
$\longrightarrow$ K_5' 

$\xrightarrow{\overline{K_5'}}$ 

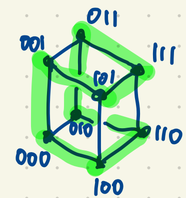
$|\mathrm{Aut}\ \overline{K_5'}| = 120$

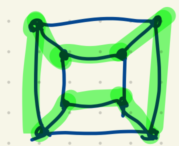
$K_5 =$ 

$\overline{K_5'} =$ 



The Petersen graph P has a Hamilton path $(0,1,2,3,6,8,5,7,9)$ (a path touching each vertex exactly once) but no Hamilton circuit (ending at the same vertex where it started).

The Hamming cube H_3 = 

= 

does have a Hamilton circuit.

000
100
110
010
011
111
101
001
000

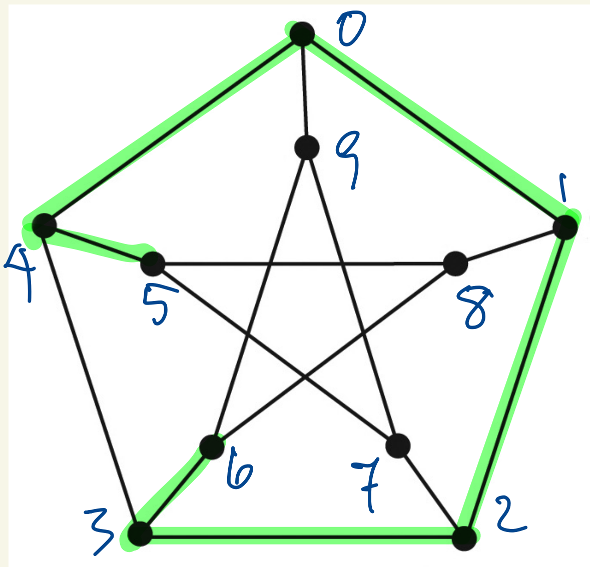
"Gray code"

Every Hamming graph H_n $(n \geqslant 2)$ has a Hamilton circuit.

A graph having a Hamilton circuit is called <u>Hamiltonian</u>.

Looking for Hamilton paths or circuits is known to be difficult in general.

Testing whether a giving graph Γ is Hamiltonian is NP-complete.

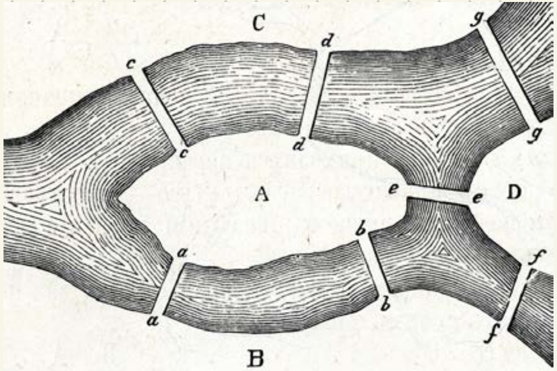


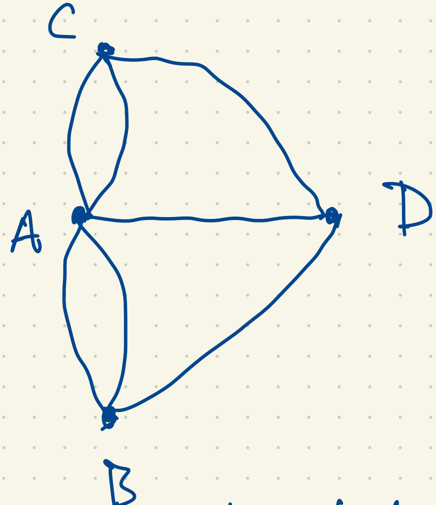
<u>Theorem</u>: The Petersen graph P is not Hamiltonian, i.e. it does not have a Hamilton circuit/cycle.

<u>Proof</u> Suppose P has a Hamilton circuit. Without loss of generality this circuit contains the path (4, 0, 1, 2). (This is because P has 120 automorphisms mapping any such path of length 3 to any other.) The Hamilton circuit uses two of the edges from vertex 3, so it uses either {3,4} or {2,3}; so without loss of generality, it uses the edge {2,3}. We cannot use the edge {3,4} as this would complete the circuit without passing through all vertices; So we must use the edges {3,6} and {4,5}. To continue the circuit from vertex 6, we have two choices: proceed through vertex 8 or vertex 9. Neither of these choices leads to a Hamilton circuit. This is a contradiction. □

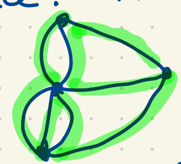
Euler paths and circuits

The Seven bridges of Königsberg





An **Euler trail** is a trail (repeating vertices but not edges) which uses each edge exactly once. An Euler circuit is an Euler trail that returns to its starting point.



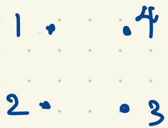
This graph has an Euler trail. In order to have an Euler trail, a graph must have either 0 or 2 vertices of odd degree. When there are no vertices of odd degree, we have an Euler circuit.

Theorem (Euler) A graph has an Euler trail iff it is connected and it has either 0 or 2 vertices of odd degree. In the case every vertex has even degree, we have an Euler circuit/cycle.

We sometimes speak of labelled graphs and unlabelled graphs.

Eg. on the vertex set $\{1,2,3,4\}$, there are $2^6 = 64$ labelled graphs

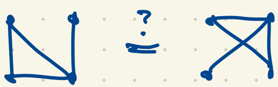
$\binom{4}{2} = 6$ pairs of vertices.



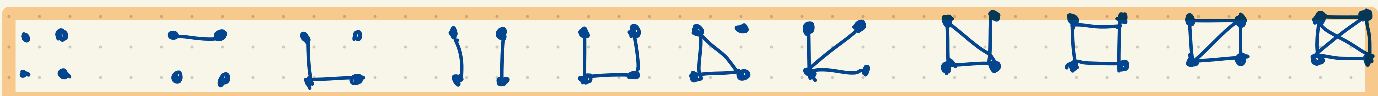
There are $\binom{n}{2}$ labelled graphs on n vertices.

But many of them are isomorphic.

These are different labelled graphs but they are isomorphic.



As unlabelled graphs they are isomorphic, hence the same graph.



There are 11 unlabelled graphs of order 4 i.e. 11 isomorphism types of graphs of order 4, i.e. 11 graphs of order 4 up to isomorphism.